

Reliable Dynamic Embedded Data Processing Systems

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Joint work with

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and others

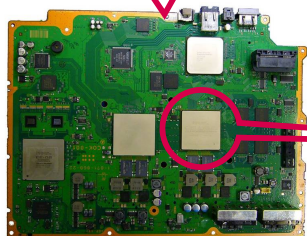
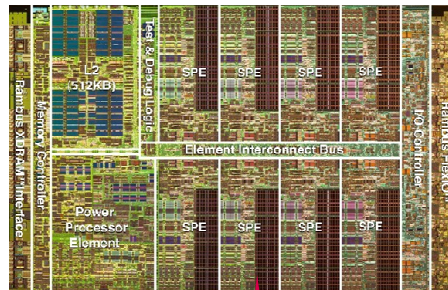
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NWO PROMES
EC FP6 Betsy, FP7 MNEMEE
SenterNovem Octopus

'Knowing is not understanding.'
Charles Kettering

2 Embedded Data Processing Systems



3 Embedded Data Processing Systems

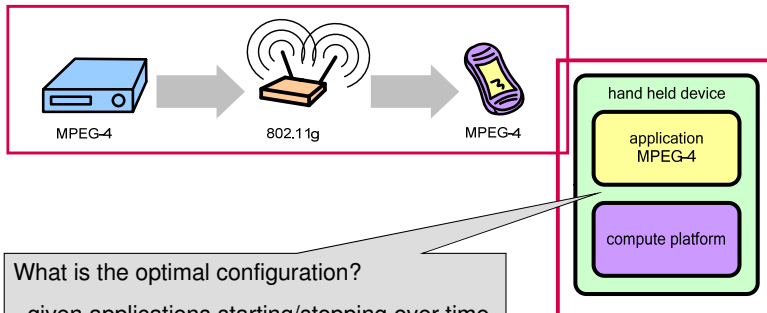


4 Embedded Data Processing Systems



5 Trend: Dynamic Behavior

- Multiprocessor systems / networking / concurrency
- Application convergence / application evolution



What is the optimal configuration?

- given applications starting/stopping over time
- given data-dependent execution times
- given varying communication bandwidth
- given battery status
- ...

- Intra-application dynamism
- Inter-application dynamism

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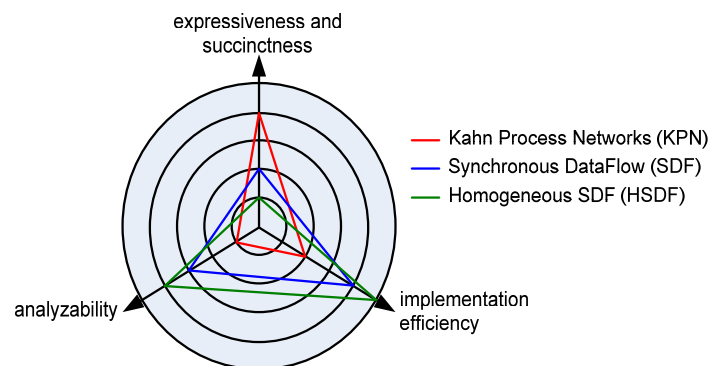
Intra-application dynamism

Dataflow analysis

7 Dataflow Analysis: Challenges

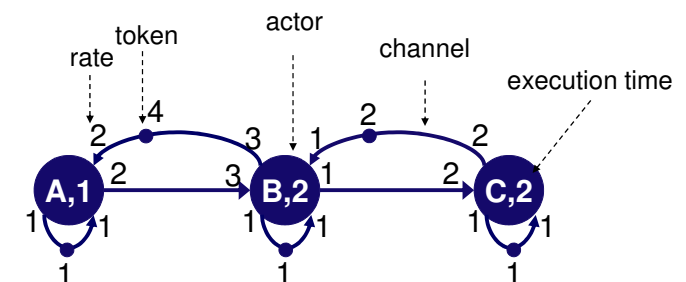
predictability and efficiency

... are contradictory

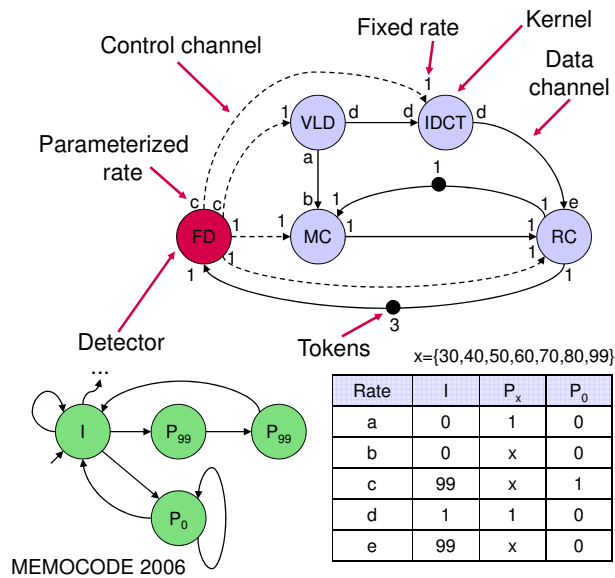


8 Synchronous Data Flow Graphs

(SDFGs [Lee 1986])



(FSM-based Scenario-Aware DataFlow (SADF))



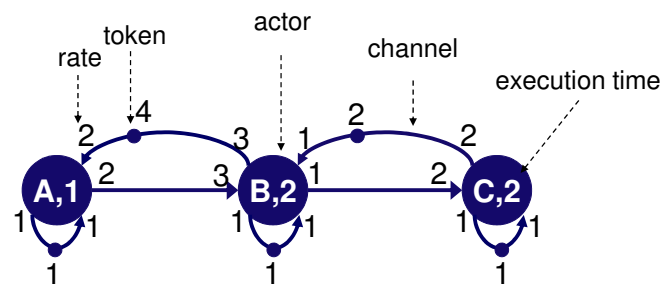
Execution time			
VLD	P_0	0	
	others	40	
IDCT	P_0	0	
	others	17	
MC	I, P_0	0	
	P_{30}	90	
	P_{40}	145	
	P_{50}	190	
	P_{60}	235	
	P_{70}	265	
	P_{80}	310	
	P_{99}	390	
RC	I	350	
	P_0	0	
	P_{30}, P_{40}, P_{50}	250	
	P_{60}	300	
FD	P_{70}, P_{80}, P_{99}	320	
	All	0	

Throughput analysis for SDF

ACSD 2006

11 Synchronous Data Flow Graphs

(SDFGs [Lee 1986])

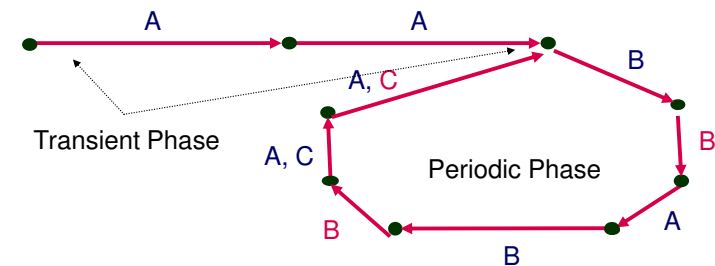


Throughput: average number of actor firings over time

Iteration: smallest non-empty set of actor firings that does not change the token distribution (example: A:3, B:2, C:1)

12 SDF Throughput Analysis

Self-timed execution:



Throughput can be calculated from the periodic phase

Efficient implementation

Consider one designated firing per iteration only to detect a recurrent state

13 Throughput Analysis: Our Result

		Traditional methods		
	State Space	Dasdan Gupta	Howard	Young Tarjan Orlin
MP3 dec.	1.10^{-3}	1.10^{-3}	1.10^{-3}	1.10^{-3}
Modem	1.10^{-3}	82.10^{-3}	81.10^{-3}	81.10^{-3}
Sample Rate	2.10^{-3}	>1800	>1800	>1800
Satellite	56.10^{-3}	>1800	>1800	>1800
H.263 decoder	10.10^{-3}	>1800	>1800	>1800

Runtimes of various throughput analysis methods (in seconds)

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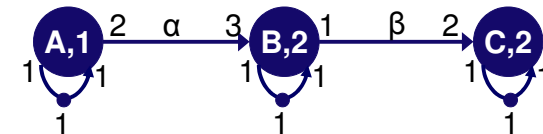
Storage-throughput trade-offs

IEEE T Comp 2008

15 Lessons to be Learned

- Repeated application of throughput analysis
- Causal dependencies potentially reducing performance

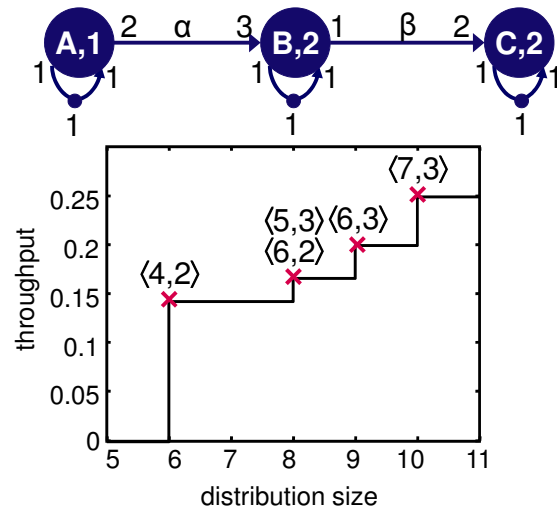
16 Trade-off: Storage vs. Throughput



Tokens on channels must be stored in memory.

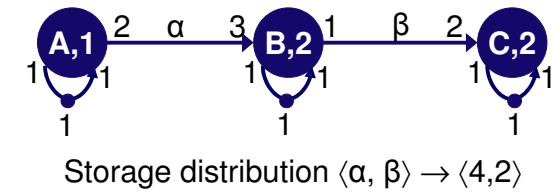
- Separate memory for each channel.
- Storage distribution, for example, $\langle \alpha, \beta \rangle \rightarrow \langle 4, 2 \rangle$

17 Trade-off: Storage vs. Throughput



Find all minimal storage distributions for any possible throughput

18 Dependency Graph



firings in one period

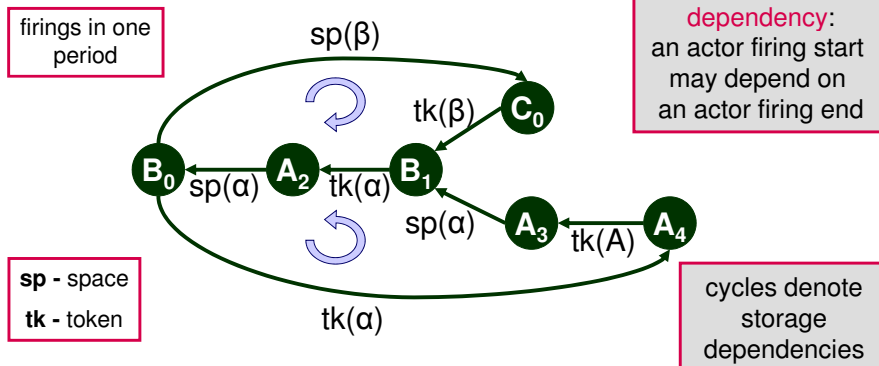
dependency:
an actor firing start
may depend on
an actor firing end

sp - space
tk - token

cycles denote
storage
dependencies

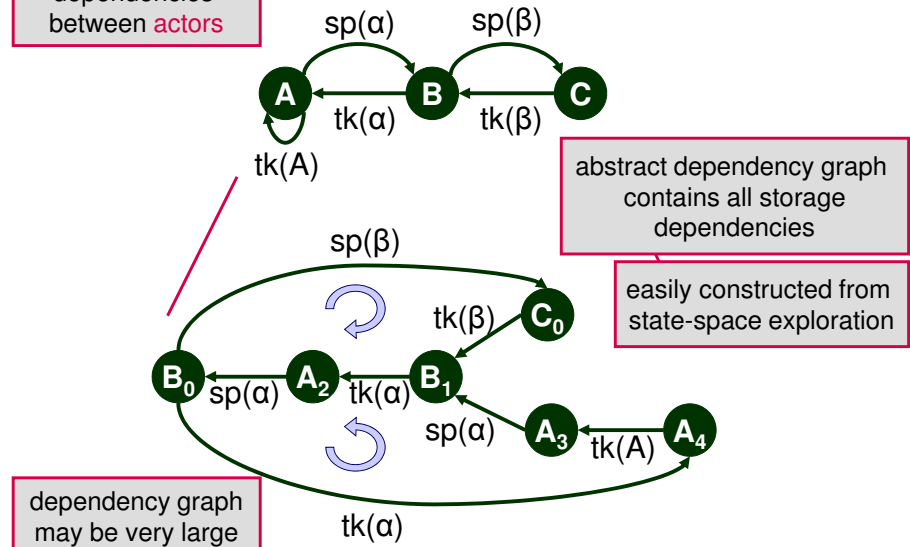
19 Dependency Graph

resolving storage dependencies may increase throughput



20 Abstract Dependency Graph

dependencies
between actors



abstract dependency graph
contains all storage
dependencies

easily constructed from
state-space exploration

dependency graph
may be very large

21 Experimental Results

	Example	Satellite	MP3	H.263
#actors	3	22	13	4
#channels	2	26	12	3
min throughput > 0	1/7	0.18	1/137	1/21876
Distr. size	6	1542	12	4753
Max through	Approximation increase buffers with n times the minimal step size			1/10000
Distr. size				8006
#Pareto points				3254
#Distributions	5	2	3	3254
Exec. time	1ms	7ms	2ms	53min

22 Trade-off Analysis

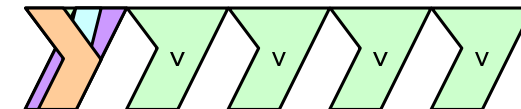
- Start small, recursively resolve bottlenecks
- Fast in general, approximation possible
- Generalizable to other resources, computational models

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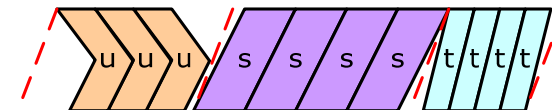
Scenarios

24 Scenario-aware Throughput Analysis

- SDF throughput analysis considers worst-case actor execution times

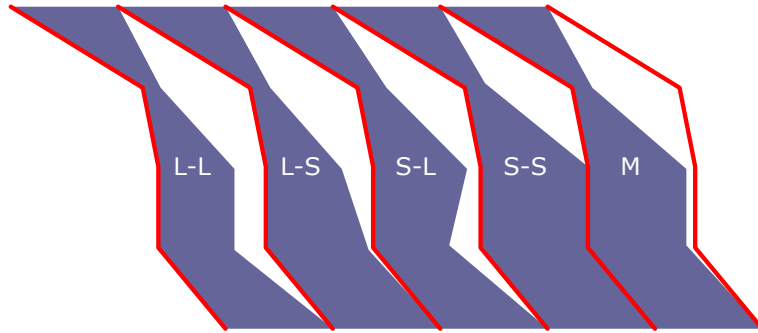


- Scenario-aware throughput analysis considers worst-case execution times per scenario, but it must consider scenario transitions

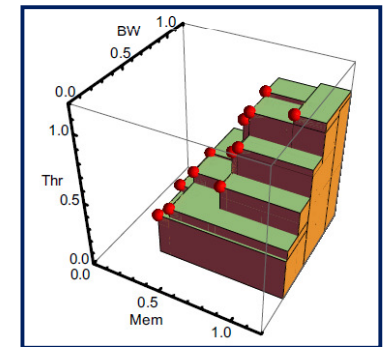


- Analyzing all scenario transitions separately can be avoided
- Separate scenario transitions with an invariant reference schedule

- SDF model: 9.28 Mcycles per frame
- SADF model: 5.83 Mcycles per frame



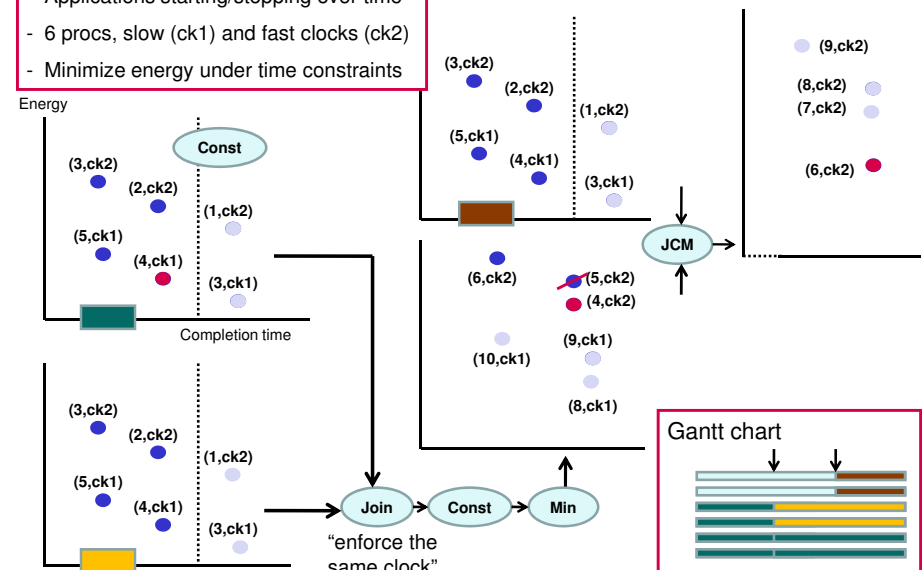
- Yang Yang: shared resources, trade-off analysis
- Sander Stuijk: SDF3 toolset, <http://www.es.ele.tue.nl/sdf3>
- Future: shared resources + scenarios + trade-off analysis



Inter-application dynamism

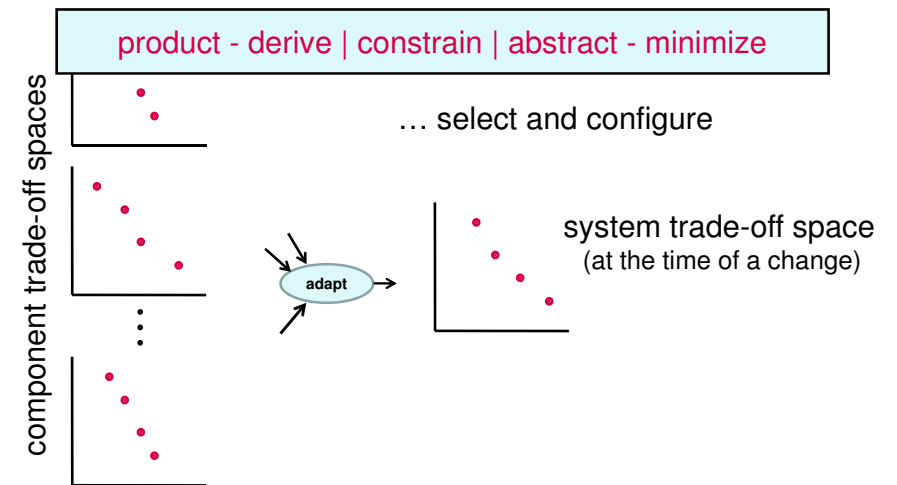
Run-time adaptation

- Applications starting/stopping over time
- 6 procs, slow (ck1) and fast clocks (ck2)
- Minimize energy under time constraints

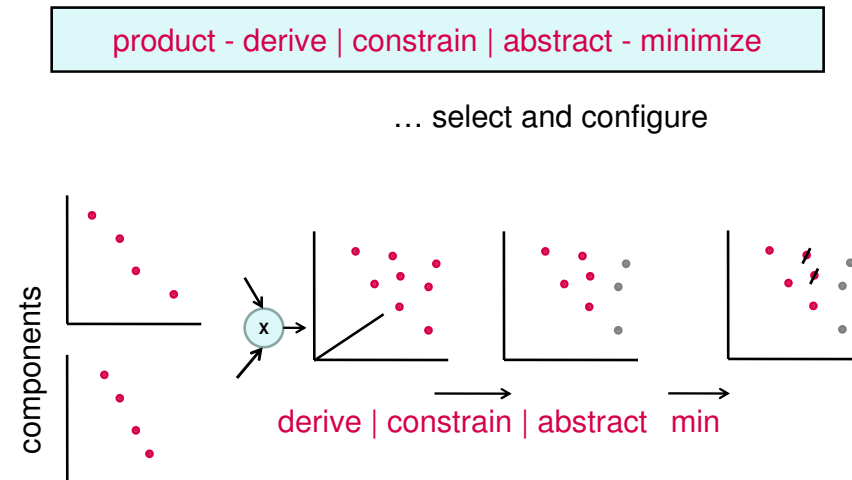


Goal: Compositional computation of trade-offs

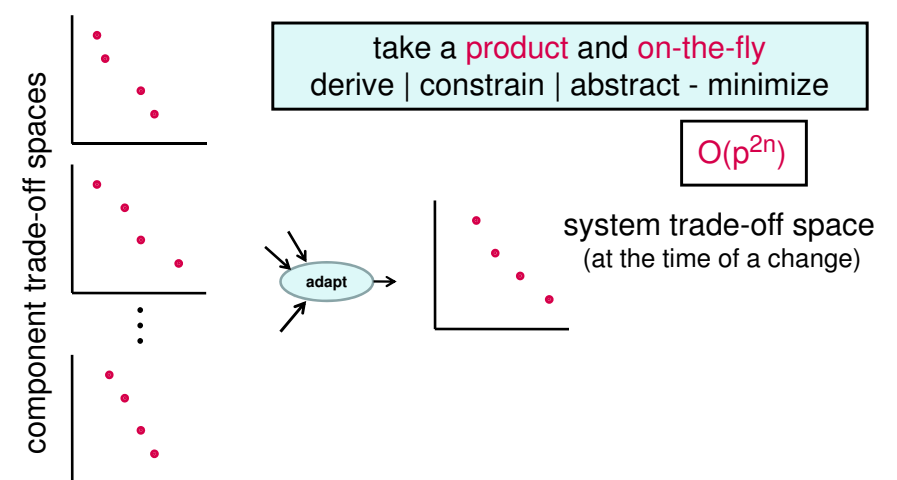
... of run-time adaptation



... of run-time adaptation



n components with p Pareto points each

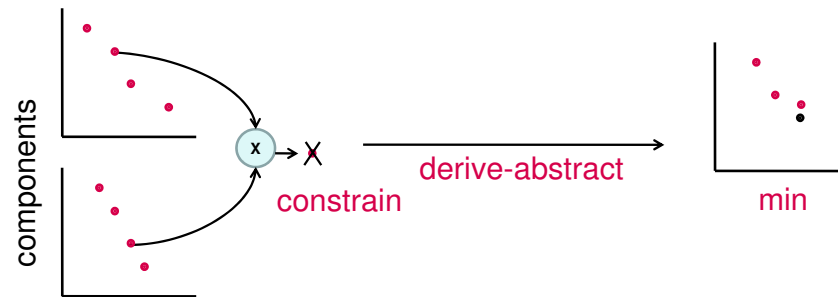


33 Complexity

n components with p Pareto points each

take a **product** and **on-the-fly**
derive | constrain | abstract - minimize

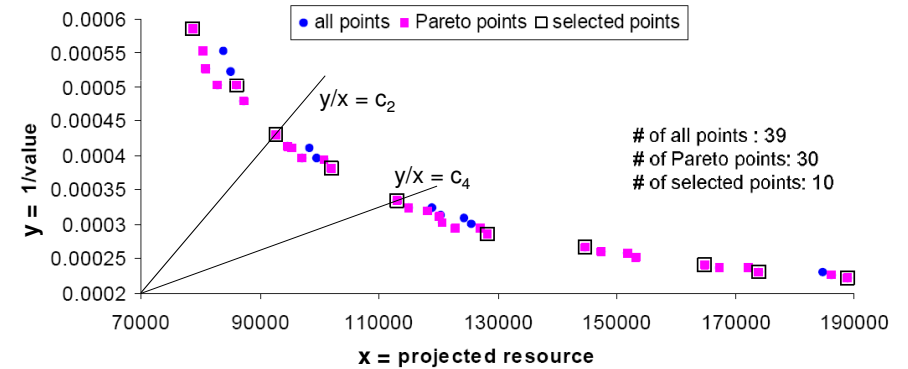
$O(p^{2n})$



34 Complexity Control

keep n and p small

- Compositional reasoning, among others
- Approximate Pareto sets



35 Multi-dimensional Multiple-choice Knapsack Problem

MMKP

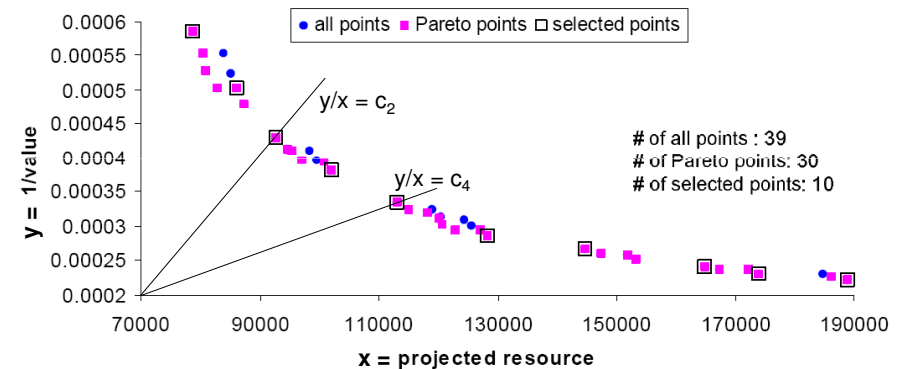
- One optimization objective (**value**)
- Multiple resource dimensions with capacity **constraints**
- Multiple independent applications
- Multiple independent configurations per application

Pick one configuration per application optimizing value within constraints

NP hard

36 A Parameterized Compositional Heuristic

- Project all resource dimensions into one dimension
- Approximate Pareto set in 2-dimensional space



product-constrain-derive-abstract-minimize

37 A Parameterized Compositional Heuristic Embedded Systems INSTITUTE TU/e

- Project all resource dimensions into one dimension
- Approximate Pareto set in 2-dimensional space

parameter **p**: maintain (at most) **p** Pareto points

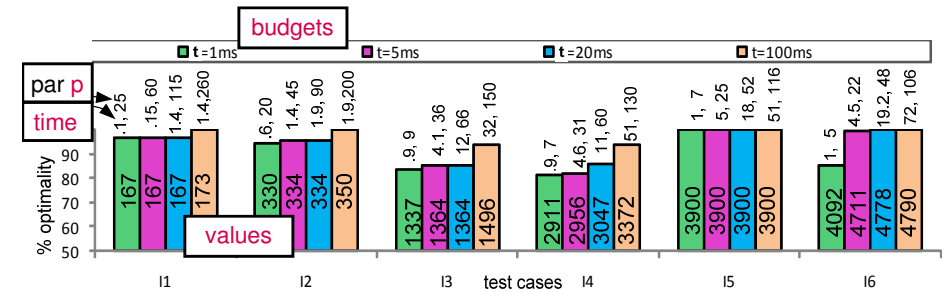
Allows to **budget** analysis time

$$\text{analysis time} \leq c \cdot n \cdot p^2$$

(*n* number of applications; *c* platform constant)

product-constrain-derive-abstract-minimize

38 Controlling Time Budgets Embedded Systems INSTITUTE TU/e

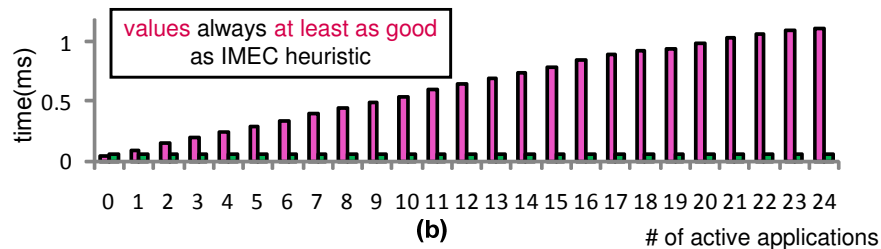
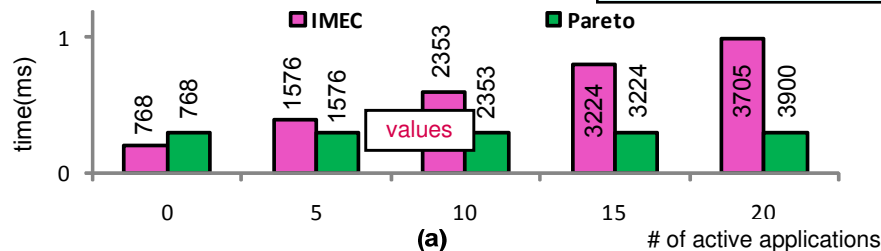


- MMKP benchmark: www.laria.u-picardie.fr/hifi/OR-Benchmark/MMKP
- Always within budget
- Good values
- Similar results to the IMEC, SOC 2006, non-compositional state-of-the-art heuristic

SimIt-ARM simulator
206 Mhz StrongARM

39 Compositionality Embedded Systems INSTITUTE TU/e

- Applications start at various times
- 5 at a time (a), 1 at a time (b)



40 Conclusions Embedded Systems INSTITUTE TU/e

- **Goal**: reliable, resource efficient data processing systems
- **Dynamic behavior** complicates analysis
- **Intra-application dynamism**: scenario-aware dataflow analysis and design-space exploration
- **Inter-application dynamism**: Pareto-algebra-based run-time adaptation

Questions ?

More info: www.es.ele.tue.nl/~tbasten/
SDF3 toolkit: www.es.ele.tue.nl/sdf3/
Pareto calculator: www.es.ele.tue.nl/pareto/

'An understanding of the natural world and what's in it
is a source of not only a great curiosity but great fulfillment.'
David Attenborough