

InvasIC seminar, Erlangen, November 2014

Making Data Flow, Dynamically



Twan Basten

Eindhoven University of Technology & TNO Embedded Systems Innovation

Joint work with
Marc Geilen, Sander Stuijk, Bart Theelen
and many others

'Knowing is not understanding.'
Charles Kettering

Electronic Systems

2

Predictable computing !?

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3 Data-intensive systems are everywhere!

Philips
medical imaging



Thales
radar



Océ
printing



Mercedes
automotive



ASML
chipfabrication



Apple
mobile computing

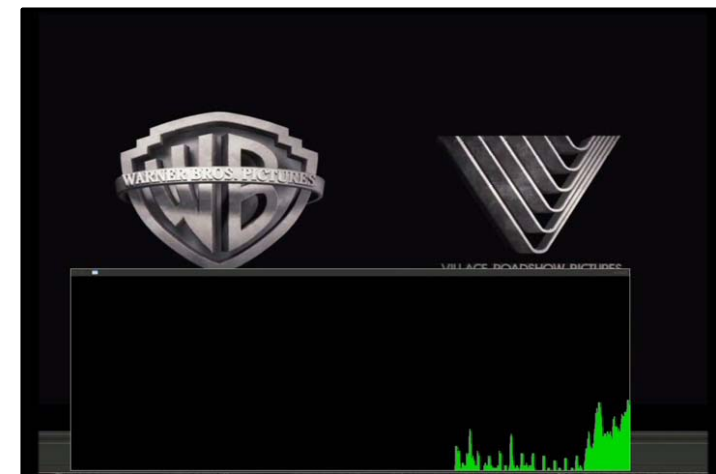


Sony
gaming



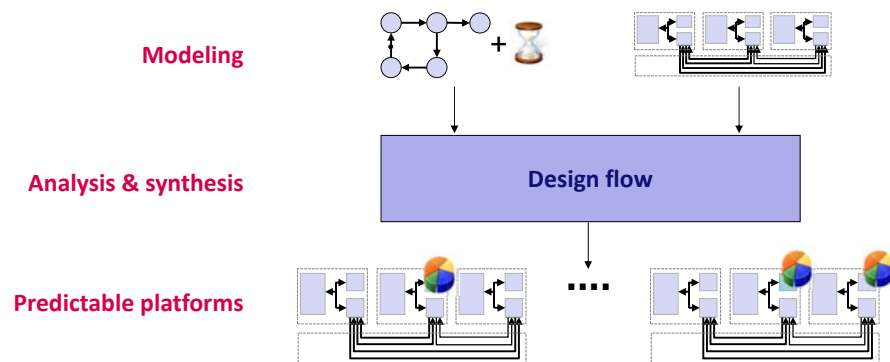
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4 Challenge: Data-dependent data-processing workloads



[Thx to Martijn Koedam]

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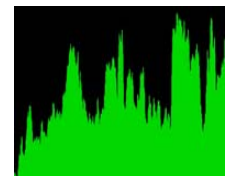
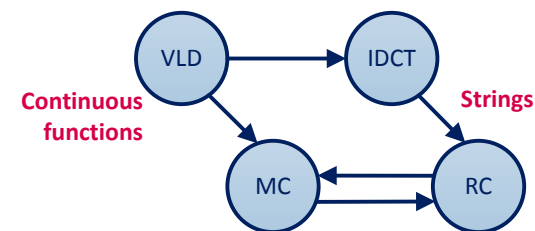


- Predictable computing !?
- Data flow models of computation
- SDF analysis and synthesis
- SADF analysis and synthesis
- Interaction
- Conclusions

Data flow models of computation

(KPN [Kahn 1974])

MPEG-4 SP:

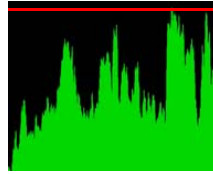
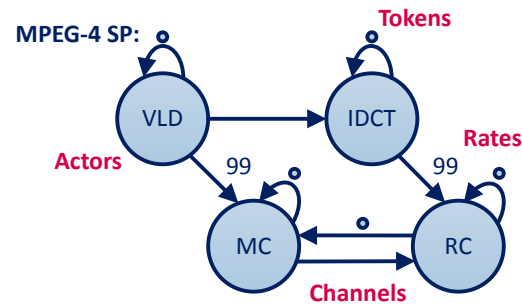


Adding **interaction**:
 - Reactive Process Networks (RPN)
 - Marc Geilen, 2004

- Can express **any continuous function** on strings
- Captures **"precisely"** all streaming data flow applications
- **Doesn't capture interaction** with the environment
- **No** analysis & synthesis (other than functionality)

9 Synchronous Data Flow

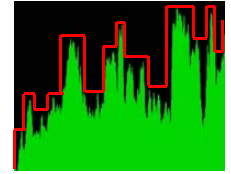
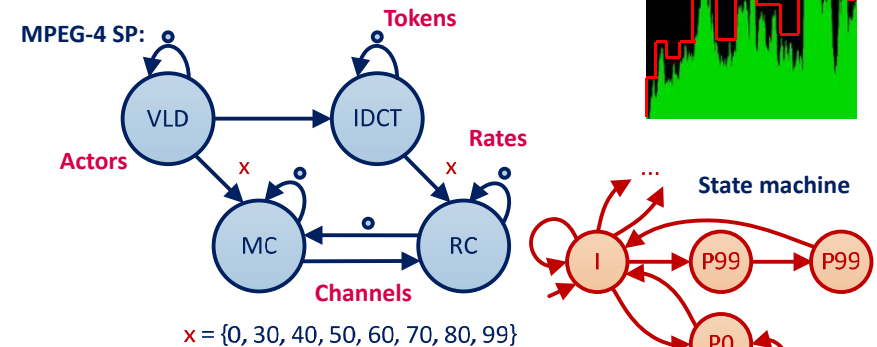
(SDF [Lee 1986])



- **Conservative**, worst-case abstraction
- **Strong analysis & synthesis** support
- Low implementation overhead ...
- ... but may lead to **over-allocation of resources**

10 Scenario-Aware Data Flow

(SADF [Theelen 2006])



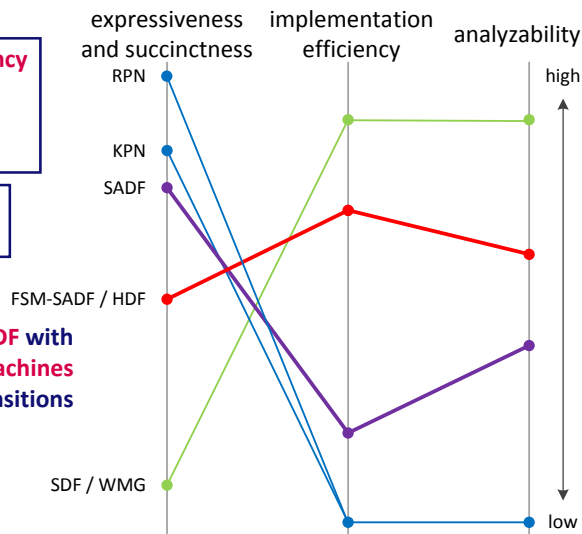
- **Dynamics** captured in **scenarios**
- **Good compromise** between expressivity, analysis & synthesis
- **Efficient implementations**

11 Comparison of data flow Models of Computation

implementation efficiency and analyzability vs. expressiveness are contradictory

but FSM-SADF provides a good compromise

SADF with Finite State Machines specifying scenario transitions



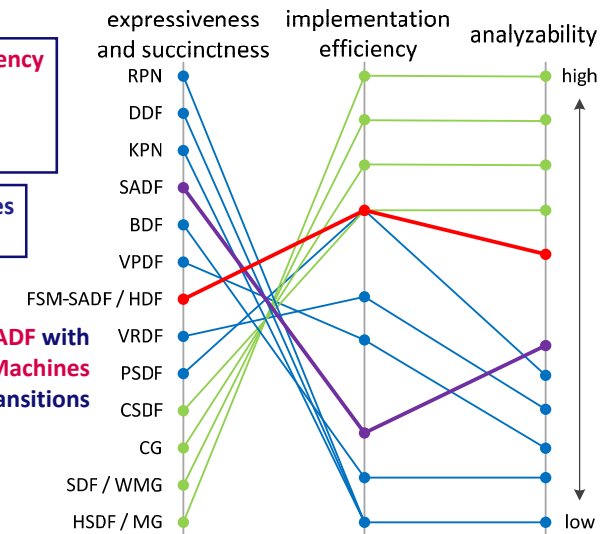
(SAMOS 2011)

12 Comparison of data flow Models of Computation

implementation efficiency and analyzability vs. expressiveness are contradictory

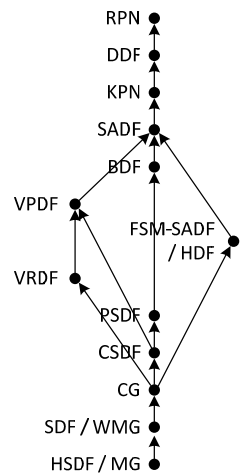
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SADF with Finite State Machines specifying scenario transitions



(SAMOS 2011)

13 Data flow Models of Computation: Expressiveness hierarchy TU/e

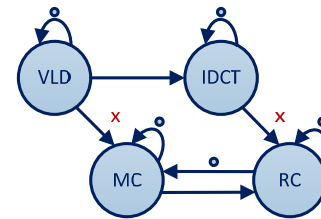


- Known or straightforward **inclusion relations**
- **BDF and above:** Turing complete
- Better expressiveness notions needed
- Unification needed

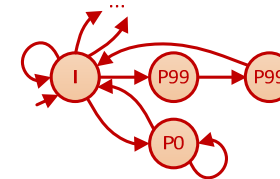
(SAMOS 2011)

14 Targeting models to the problem at hand **is important** TU/e

2-processor implementation of MPEG4-SP



$x = \{0, 30, 40, 50, 60, 70, 80, 99\}$



Actor **worst-cases**
occur for different scenarios

SDF	
Execution times	
VLD	33
IDCT	14
MC	325
RC	292

Data rate	
x	99

SADF	Execution times		
	VLD	P0	0
		other	33
	IDCT	P0	0
		other	14
	MC	I,P0	0
		P30	75
		P40	121
		P50	158
		P60	196
P70		221	
P80		258	
P99		325	
RC		I	292
	P0	0	
	P30,40,50	208	
	P60	250	
	P70,80,99	267	

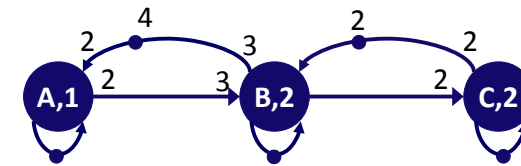
15 Targeting models to the problem at hand is important TU/e

MPEG4-SP				MP3
metric	SDF	SADF	improvement	improvement
throughput	15 (frames/s)	15 (frames/s)	NA	NA
processor capacity	88 %	26 %	70 %	-
memory capacity	254 KB	254 KB	-	21 %
bandwidth	33 %	33 %	-	23 %

16 Outline TU/e

- Predictable computing
- Data flow models of computation
- SDF analysis and synthesis
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SDF analysis and synthesis



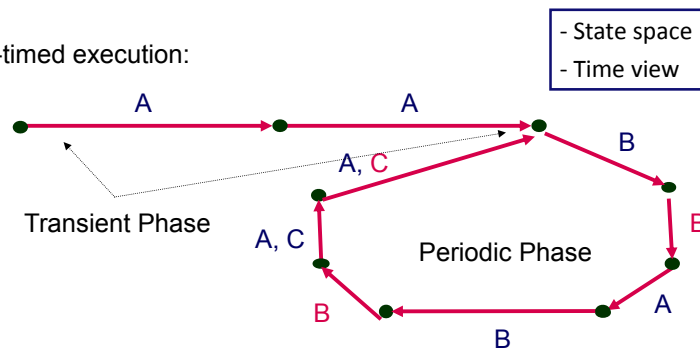
Iteration

smallest non-empty set of actor firings
that does not change the token distribution

(example: A:3, B:2, C:1)

Crucial concept in data flow analysis

Self-timed execution:

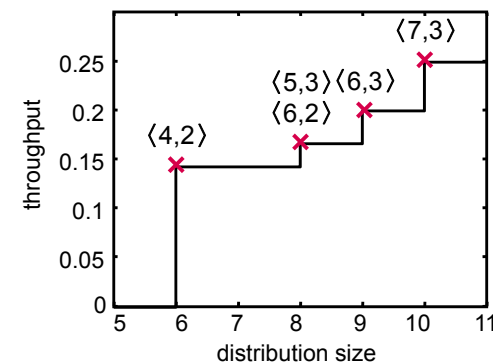


Throughput: average number of actor firings over time

Throughput can be calculated from the periodic phase

Efficient implementation

Consider one designated firing per iteration only to detect a recurrent state



- Separate buffer per channel
- Storage distribution, e.g., $\langle \alpha, \beta \rangle \rightarrow \langle 4, 2 \rangle$

Find all minimal storage distributions for any possible throughput

21 Storage vs. throughput: Dependency analysis

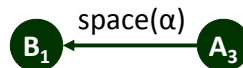
TU/e



Storage distribution $\langle \alpha, \beta \rangle \rightarrow \langle 4, 2 \rangle$

dependency:

an actor firing start may depend on an actor firing end



e.g., the 3rd firing of A, A₃, depends on the first firing of B, B₁

space limitation

resolving space dependencies may increase throughput

Efficient implementation

- capture dependencies during state-space exploration ...
- at actor level

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(DAC 2006, TC 2008)

22 Storage vs. throughput: Dependency analysis

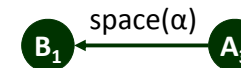
TU/e



Storage distribution $\langle \alpha, \beta \rangle \rightarrow \langle 4, 2 \rangle$

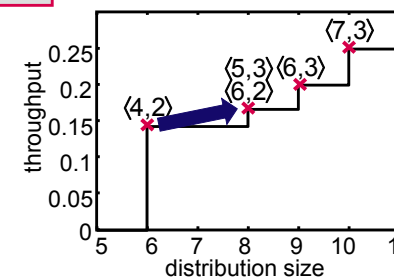
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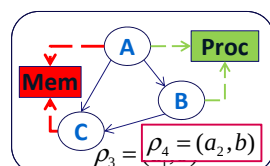
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(DAC 2006, TC 2008)

23 Trade-off analysis – resource sharing

TU/e

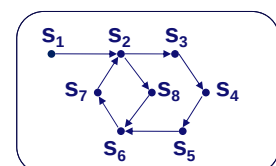
(Resource-Aware SDF, RASDF [Yang 2009])



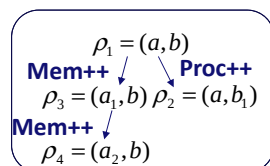
RASDF model



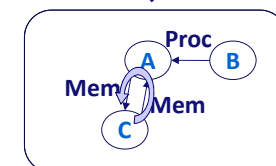
Etc



State-space exploration



Dimensioning



Dependency analysis

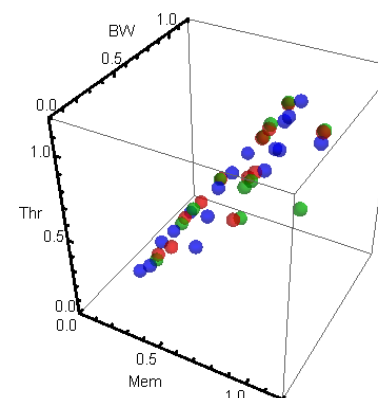
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(DATE 2010)

24 Trade-off analysis – resource sharing

TU/e

(Resource-Aware SDF, RASDF [Yang 2009])



- Printer architecture exploration
- 3 architectures (colors)
- Throughput, memory, bandwidth trade offs

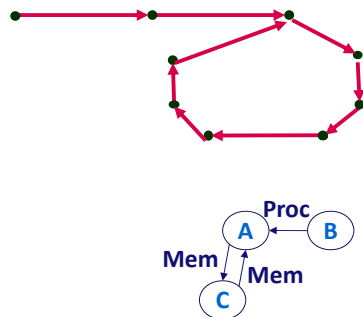
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(DATE 2010)

25 Lessons learned: key concepts

TU/e

- Iterations
- State-space exploration
- Dependency analysis



These concepts are re-usable !

26 Outline

TU/e

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27

TU/e

SADF analysis and synthesis

28 SADF throughput analysis

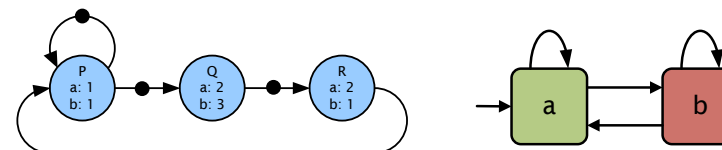
TU/e

- Analysis based on (max, +)-algebra



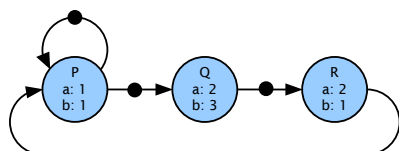
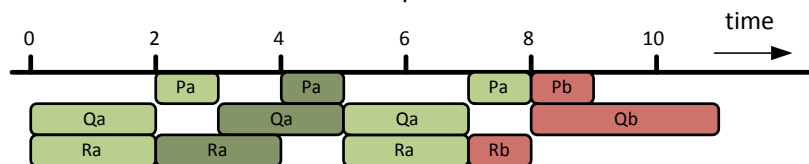
$$t_3 = t_4 = \max(t_1, t_2) + E$$

- SADF graph with two scenarios a and b
- Each iteration executes in one scenario



(CODES+ISSS 2010)

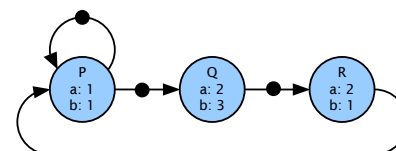
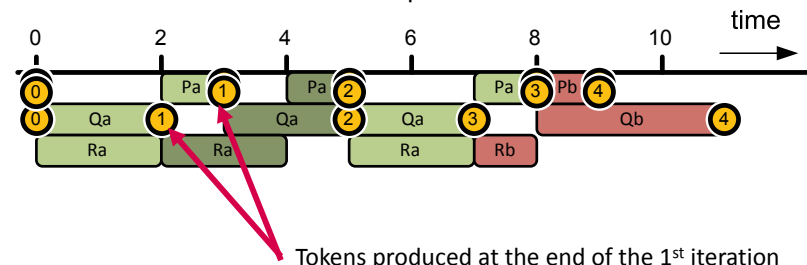
- Gantt chart for the scenario sequence **aaab**



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(CODES+ISSS 2010)

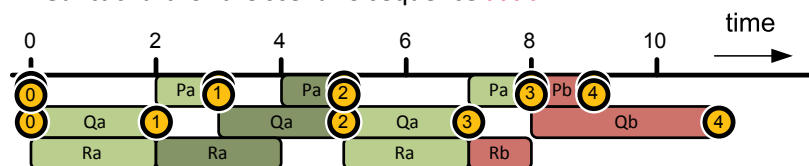
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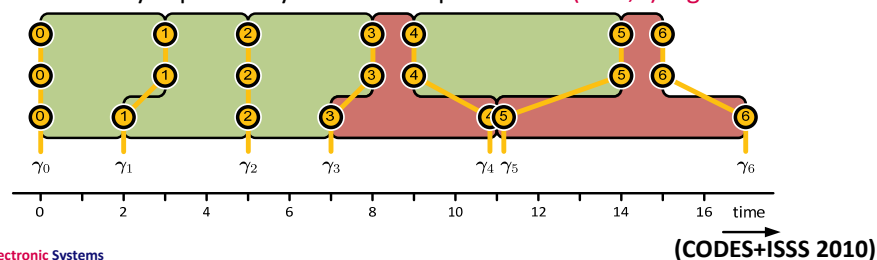
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(CODES+ISSS 2010)

- Gantt chart for the scenario sequence **aaab**



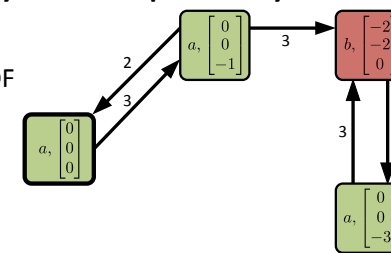
- Execution is a **sequence of vector shapes**
- Is nicely captured by matrix multiplication in **(max,+)-algebra**



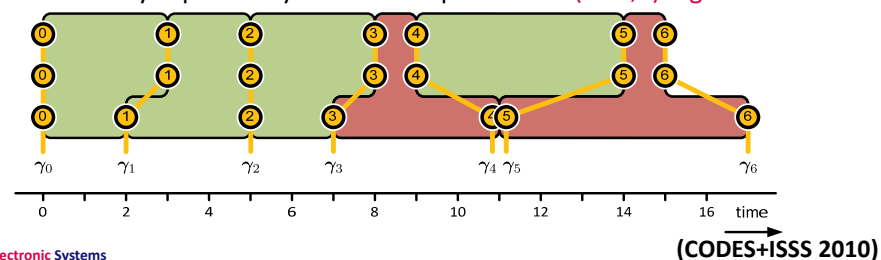
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(CODES+ISSS 2010)

- State space** of the SADF



- Execution is a **sequence of vector shapes**
- Is nicely captured by matrix multiplication in **(max,+)-algebra**



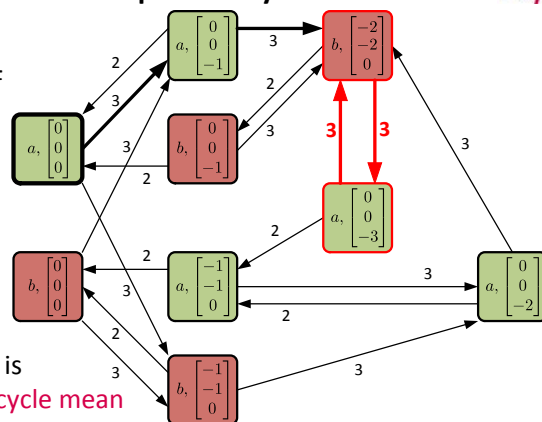
Electronic Systems

(CODES+ISSS 2010)

33 SADF throughput analysis: state-space analysis

TU/e

- State space of the SADF



- Worst-case throughput is the inverse of the max cycle mean

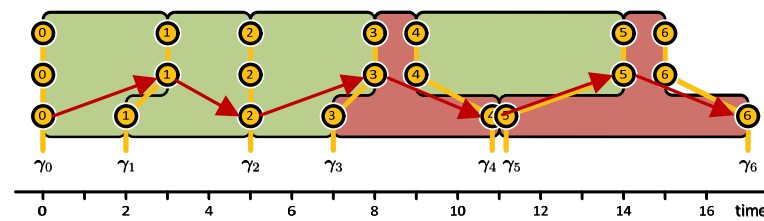
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(CODES+ISSS 2010)

34 SADF throughput analysis: (Max,+)-automata

TU/e

- Critical path is through token dependencies



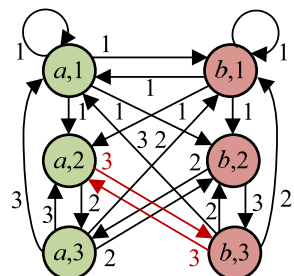
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(CODES+ISSS 2010)

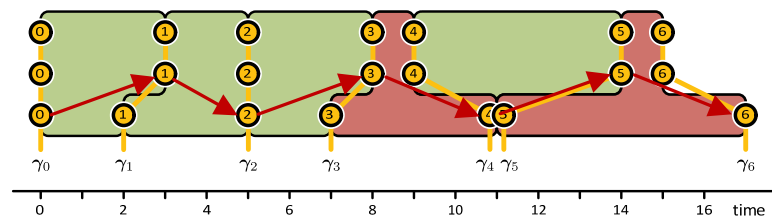
35 SADF throughput analysis: (Max,+)-automata

TU/e

- (Max,+)-automaton
- Worst-case throughput is again the inverse of the max cycle mean



- Critical path is through token dependencies



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(CODES+ISSS 2010)

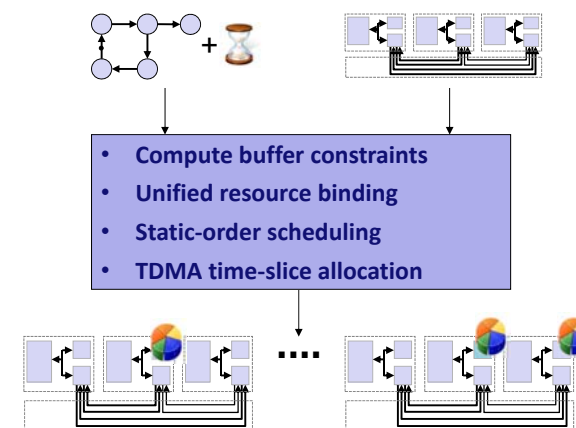
36 Model-driven SADF design flow: the SDF3 toolkit

TU/e

Modeling

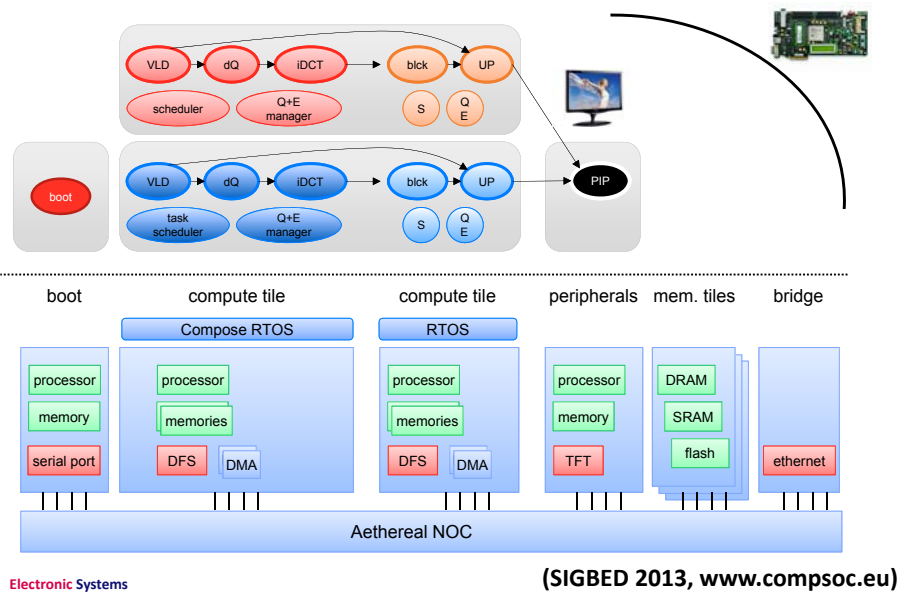
Analysis & synthesis

Predictable platforms



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(DSD 2010, www.es.ele.tue.nl/sdf3)



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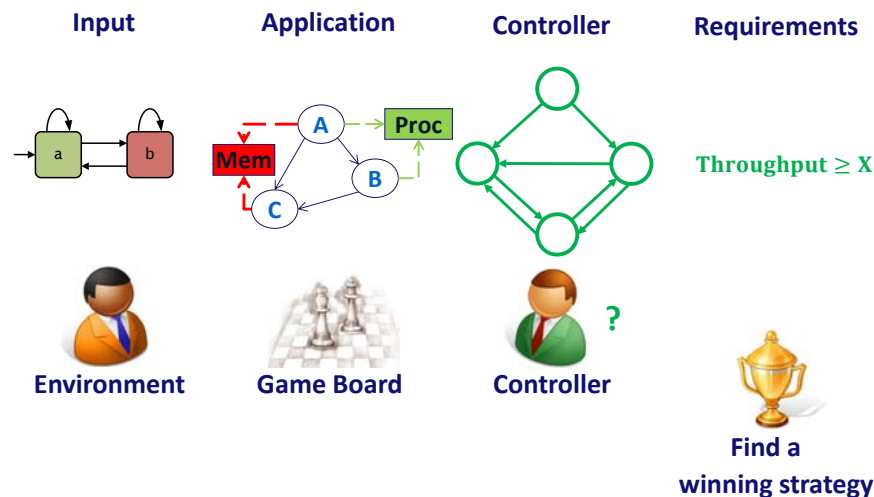
Interaction



- Can throughput be guaranteed ?
- Can energy usage be optimized under a throughput constraint ?

41 Realizing guaranteed throughput becomes a game

TU/e



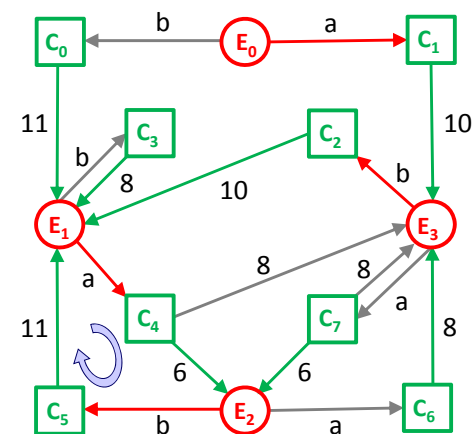
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(DATE 2012)

42 Controller is constructed by solving a mean payoff game

TU/e

- Game graph can be derived from SADF graph
- For every input scenario, controller executes a pre-defined schedule with the indicated latency
- Solved by policy iteration
 - Red encodes the worst-case input
 - Green encodes the optimal controller
- Guaranteed throughput: inverse of the maximum cycle mean



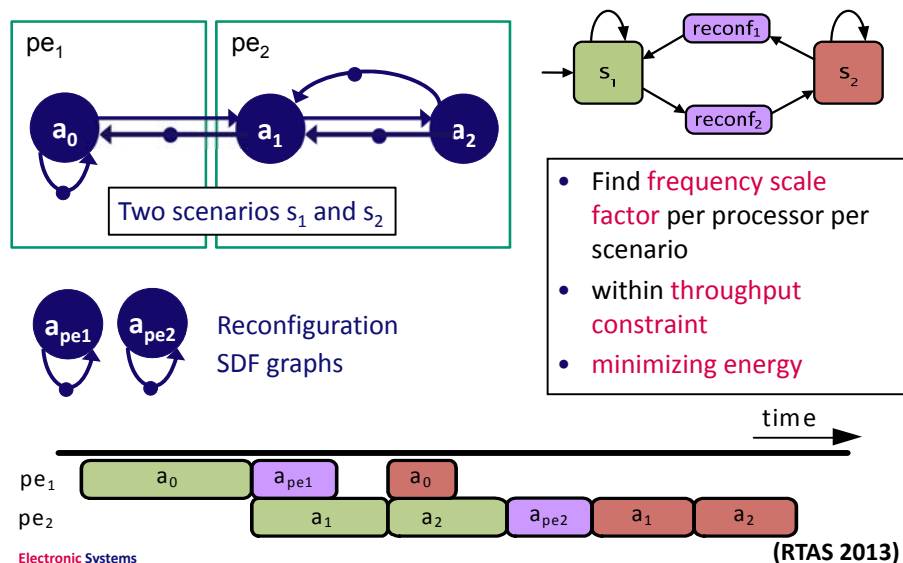
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(DATE 2012)

43 Dynamic Frequency and Voltage Scaling

TU/e

SADF model of mapped application, including reconfigurations



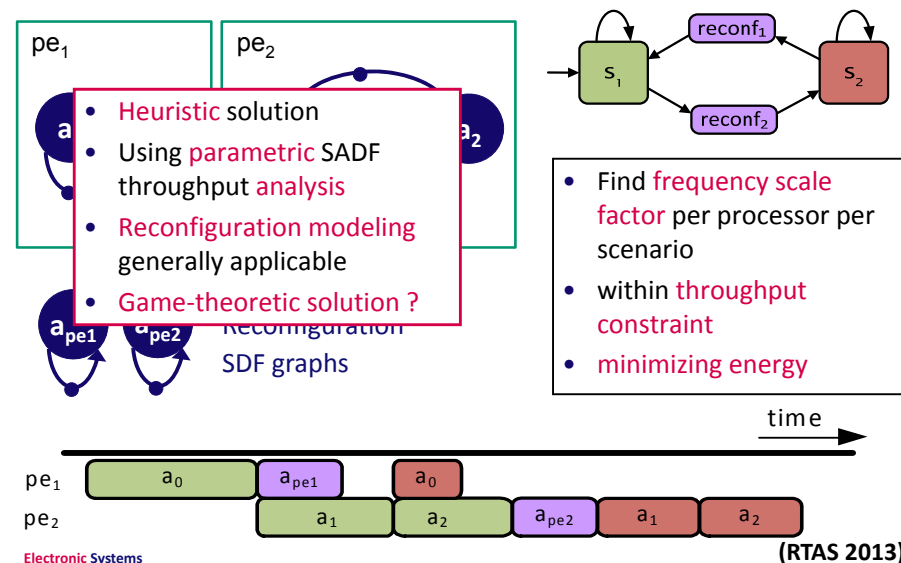
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(RTAS 2013)

44 Dynamic Frequency and Voltage Scaling

TU/e

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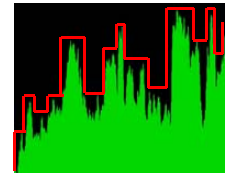
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(RTAS 2013)

- Predictable computing
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- **Data flow** MoCs:
 - Increasingly important
 - Good compromise between expressivity and analyzability
 - Synthesis feasible, leading to efficient implementations
- **SDF3**: Fully operational flow from **SDF** to **CompSOC** platform
- **Challenges**: Coping with **dynamics**
 - Expressivity, unification of MoCs
 - Modeling mapping decisions
 - Predictable reconfiguration
 - Buffer & memory analysis
 - Resource sharing
 - Controller synthesis
 - Multi-application synthesis
 - ...

It's an exciting field !



Questions ?

More info: www.es.ele.tue.nl/~tbasten/

SDF3 toolkit: www.es.ele.tue.nl/sdf3/

CompSOC: www.compsoc.eu



Twan Basten



Marc Geilen

Sander Stuijk

